

**Problem 1.(15 Points)**

The following is the algorithm for sorting by insertion

**procedure** *insertion sort* ( $a_1, a_2, \dots, a_n$ : real numbers with  $n \geq 2$ )

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1   for  $j := 2$  to  $n$ 
2   begin
3        $position := 1$ 
4       while  $a_j > a_{position}$ 
5            $position := position + 1$ 
6        $m := a_j$ 
7       for  $k := 0$  to  $j - position - 1$ 
8            $a_{j-k} := a_{j-k-1}$ 
9        $a_{position} := m$ 
10  end {  $a_1, a_2, \dots, a_n$  are sorted }
```

For each  $j$ , the algorithm does the following:  $a_1, \dots, a_{j-1}$  is already sorted, it inserts  $a_j$  in the “correct position” so that  $a_1, \dots, a_j$  is sorted. The “correct position” is computed in lines 3-4-5 in the variable  $i$ . It is “searched” for in a way based on linear search.

- a. (4) Complete lines 5 and 8 above
- b. (4) Explain why the while loop in line 4 must terminate (i.e. its condition will definitely become false.)

Eventually,  $position$  will become  $j$  and so  $a_j = a_{position}$  (worst case)

- c. (4) Use the insertion sort procedure above to sort the list 6, 2, 3, 1, 4 showing the lists obtained at the end of each step. Here  $n=5$ . Fill out the following table.

$j$	$a_1$	$a_2$	$a_3$	$a_4$	$a_5$
	6	2	3	1	4
2	2	6	3	1	4
3	2	3	6	1	4
4	1	2	3	6	4
5	1	2	3	4	6

- d. (3) For the step  $j = 3$ , what will be the value of  $i$  in line at the end of the loop of lines 4-5, and what is the stopping condition that will terminate the loop?

$$i = 2$$

*Stopping condition* :  $3 > 6$  False

**Problem 2.(15 Points)**

This problem refers to the insertion sort algorithm given in the previous problem. You are supposed to do a **worst-case** complexity analysis for insertion sort, based on the count of number of comparisons used. Assume that in the **for** iterations in lines 1 and 7, one comparison is used for each loop, to check whether the “limit has been reached”.

- a. (4) For each  $j$  how many comparisons  $C_1(j)$  will be used to execute lines 3-4-5 ?

$$C_1(j) = \text{position} + 1 \text{ (position} < j)$$

- b. (4) For each  $j$  how many comparisons  $C_2(j)$  will be used to execute lines 6-7-8-9 ?

$$C_2(j) = j - \text{position} + 1$$

- c. (4) What is the overall count  $C(n)$  of comparisons used to sort a list with  $n$  elements? Show your computations.

$$\begin{aligned} C(n) &= \sum_{j=2}^n (C_1(j) + C_2(j) + 1) + 1 \\ &= \sum_{j=2}^n (\text{position} + j - \text{position} + 2) + 1 \end{aligned}$$

$$\begin{aligned} C(n) &= \frac{n(n+1)}{2} - 1 + 2(n-1) + 1 \\ &= \frac{n^2}{2} + \frac{5n}{2} - 2 \end{aligned}$$

- d. (3) Based on your answer in part (c), what is the big-O estimate that you can give for the algorithm insertion sort ? { Use **only one** of  $\log n$ ,  $n$ ,  $n \log n$ ,  $n^2$ ,  $n^3$  }

$$C(n) = O( n^2 )$$

**Problem 3.(10 Points)**

The following algorithm is based on binary search for determining the correct position in which to insert a new element  $x$  in an already sorted list.

**procedure** *binary search position* ( $x$  : integer,  $a_1, a_2, \dots, a_n$ : increasing integers)

$i := 1$  { $i$  is left endpoint of search interval}

$j := n$  { $j$  is right endpoint of search interval}

**while**  $i < j$

**begin**

$m := \lfloor (i+j)/2 \rfloor$

**if**  $x > a_m$  **then**  $i := m+1$  **else**  $j := m$

**end**

**if**  $x > a_i$  **then**  $position := i + 1$

**else**  $position := i$

{ $position$  is the correct position to insert  $x$ }

- (2) Complete the last **else** line of the algorithm above.
- (4) What is the count of comparisons for this algorithm? Why?

Let  $k$  be such that  $2^k \leq n < 2^{k+1}$  i.e.  $k = \lfloor \log_2 n \rfloor$

Number of comparisons =  $2k + 1$  (for the while loop)  
 + 1 (for the if)  
 =  $2k + 2 = 2\log_2 n + 2$

- (4) The algorithm above will be used in the insertion sort algorithm above for determining the correct position in which to insert  $a_j$  in the list  $a_1, a_2, \dots, a_{j-1}$ . So lines 3-4-5 in *insertion sort*, will be replaced by a “call” to the binary search position procedure as follows:

*binary search position* ( $a_j, a_1, a_2, \dots, a_{j-1}$ )

How would your answers to parts (c) and (d) of the previous problem change?

$$C(n) = [ \sum_{j=2}^n (2\log_2(j-1) + 2 + j - position + 1) ] + 1$$

position = 1 worst case

$$C(n) = ( 2 \sum_{j=2}^n \log_2(j-1) + 2 + j ) + 1$$

$$C(n) = O( n^2 ) \quad \{ \text{Use only one of } \log n, n, n \log n, n^2, n^3 \}$$

**Problem 4.(5 Points)**

- a. Use the definition of big-oh to prove that  $1.2 + 2.3 + \dots + (n-1).n$  is  $O(n^3)$ .

$$\begin{aligned} 1.2 + 2.3 + \dots + (n-1).n &\leq (n-1).n + (n-1).n + \dots + (n-1).n \\ &\text{[n-1 times]} \\ &= (n-1)^2.n \leq n^3 \text{ for all } n \geq 2 \end{aligned}$$

Which is  $O(n^3)$ ; witnesses  $c = 1$  and  $n \geq 2$

**Problem 5. (15 Points)**

Consider the proposition

$P(n)$ : An amount of postage of  $n$  cents can be formed using 3-cent and 7-cent stamps.

You should prove that  $P(n)$  is true for  $n \geq 12$ , first using mathematical induction and then using strong mathematical induction.

- a. (8) Use mathematical induction to prove that  $P(n)$  is true for  $n \geq 12$ .

$$P(12) = 4 \cdot 3 + 0 \cdot 7 \text{ true}$$

Assume it is true for  $k \geq 12$ . Show it for  $P(k+1)$

$$k = a \cdot 3 + b \cdot 7$$

$$\begin{aligned} \text{if } a \geq 2, \text{ then } k + 1 &= a \cdot 3 + b \cdot 7 + 1 \\ &= (a-2) \cdot 3 + (b+1) \cdot 7 \end{aligned}$$

$$\begin{aligned} \text{if } a = 0, \text{ then } k &= 7b \text{ and since } k \geq 12 \Rightarrow b \geq 2 \\ k + 1 &= 7 \cdot (b-2) + 15 = 5 \cdot 3 + (b-2) \cdot 7 \end{aligned}$$

$$\begin{aligned} \text{if } a = 1, \text{ then } k &= 3 + 7 \cdot b \text{ and again since } k \geq 12, \text{ then } b \geq 2 \\ \text{so } k + 1 &= 6 \cdot 3 + (b-2) \cdot 7 \end{aligned}$$

therefore  $P(n)$  is true for all  $n \geq 12$ .

- b. (7) Use strong mathematical induction to prove the result. (Hint: Show that the statements  $P(12)$ ,  $P(13)$ , and  $P(14)$ , are true, and use strong induction accordingly)

$$P(12) = 4 \cdot 3 + 0 \cdot 7, \text{ true}$$

$$P(13) = 2 \cdot 3 + 1 \cdot 7, \text{ true}$$

$$P(14) = 0 \cdot 3 + 2 \cdot 7$$

Assume  $p(k)$  is true for  $k = 12, 13, 14, \dots, n ; n \geq 14$

Consider  $n+1$

$$n+1 - 3 \geq 14 + 1 - 3 \geq 12$$

so  $P(n+1-3)$  is true

$$n + 1 - 3 = a \cdot 3 + b \cdot 7$$

$$n + 1 = (a+1) \cdot 3 + b \cdot 7, \text{ then } P(n+1) \text{ is true}$$

therefore  $P(n)$  is true for all  $n \geq 12$

**Problem 6. (10 Points)**

In the questions below give a recursive definition with initial condition(s).

- a. (4) The function  $f(n) = 2^n$ ,  $n = 1, 2, 3, \dots$

Basis step:  $f(1) = 2$

Recursive step:  $f(n) = 2 * f(n-1)$ ,  $n \geq 2$

- b. (3) The sequence  $a_1 = 16, a_2 = 13, a_3 = 10, a_4 = 7, \dots$

Basis step:  $a_1 = 16$

Recursive step:  $a_n = a_{n-1} - 3$ ,  $n \geq 2$

- c. (3) The set  $\{0, 3, 6, 9, \dots\}$ .

Basis step:  $0 \in S$

Recursive step:  $x \in S \rightarrow x + 3 \in S$

**Problem 7.(15 Points)**

In the questions below suppose that a “word” is any string of seven letters of the English alphabet, with repeated letters allowed. There are 26 letters in the English alphabet, and there are 5 vowels (a, e, i, o, u.) Show your computations

- a. (2) How many words are there?

$$26^7$$

- b. (2) How many words begin with R **and** end with T?

$$26^5$$

- c. (3) How many words begin with A **or** end with B?

Begin with ‘A’: 266

End with ‘B’: 266

Begin with ‘A’ and end with ‘B’: 265

Begin with ‘A’ or end with ‘B’:  $2 * 266 - 265$

- d. (4) How many words begin with the first three letters are A, A, B in some (any) order?

A, A, B in some order: 3 possibilities: AAB, ABA, BAA

So  $3 * 264$

- e. (4) How many words **have exactly** 2 vowels?

Number of possibilities for the position of these 2 vowels =

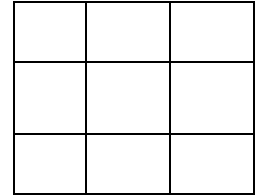
$$\binom{7}{2} = \frac{7!}{2!5!} = \frac{7!}{2!5!} = \frac{7*6}{2!} = 21$$

For each possibility: 215

So total:  $21 * 25 * 215$

**Problem 8 (15 Points)**

- a. (5) Show that if 10 points are picked on or in the interior of a square of side length 3, then there are at least two of these points no farther than  $\sqrt{2}$  apart.



The square can be divided into 9 squares each of dimensions 1 x 1.

2 of the 10 points have to be in one of these squares (by pigeon hole principle)

The distance between these two points is at most  $\sqrt{2}$  (diagonal distance – max in square)

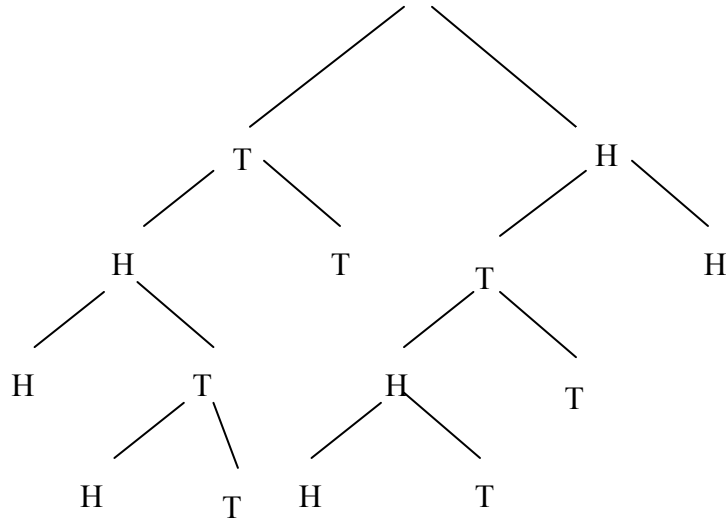
- b. (5) Find the least number of cables needed to connect 10 computers to 4 printers to guarantee that 6 computers can directly access 4 different printers. Justify your answer.

If 30 cables are given, then they can be used so that each of the 10 computers get connected to 3 printers each.

To guarantee that 6 computers are connected to 4 different printers each: 36 cables



- c. (5) A game consisting of flipping a coin ends when the player gets two heads in a row, two tails in a row, or flips the coin four times.
- (a) Draw a tree diagram to show the ways in which the game can end.
- (b) In how many ways can the game end?



8 ways

**Problem 9. (10 Points)**

- a. (5) How many functions are there from the set  $\{1, 2, \dots, n\}$  to the set  $\{1, 2, 3\}$ . Why?

For each of  $1, 2, \dots, n$  there are 3 possible images  $\Rightarrow 3^n$

- b. (5) Out of the functions in (a), how many are one-to-one? Why?

If  $n > 3$ , then 0 functions

If  $n \leq 3$ , then:

$n = 3$ :

for 1 we have 3 choices

for 2 we have 2 choices

for 3 we have 1 choice

$$3 * 2 * 1 = 6 \text{ choices}$$

$n = 2$ :  $3 * 2$

$n = 1$ : 3

$n$	$3 \dots (3 - n + 1)$
3	$3 * 2 * 1$
2	$3 * 2$
1	3